

## IMPLICIT FLOOD ROUTING IN NATURAL CHANNELS<sup>a</sup>

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Discussion by Danny L. Fread

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DANNY L. FREAD,<sup>b</sup> M. ASCE.—The authors state that the  $2N \times 2N$  coefficient matrix associated with Eq. 18 has a maximum of only four non-zero el-

<sup>a</sup> December, 1970, by Michael Ameln and Ching S. Fang (Proc. Paper 7773).

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elements in any one row and that these elements are banded around the main diagonal. This feature enables a fast matrix solution technique to be devised.

The writer has encountered this type of matrix previously and developed a direct solution technique similar to the Gauss elimination method. The technique offers the following advantages:

1. The computations do not involve any of the many zero elements in the coefficient matrix; this saves considerable computation time, and
2. The required computer core storage is reduced significantly from that required for a  $2N \times 2N$  matrix to that required for a  $2N \times 4$  matrix; this results in a  $100(N - 2/N)$  percentage reduction of storage.

The authors state that the first advantage is to be sought when devising a matrix solution procedure; however, they do not mention the second advantage. The writer has found that the reduction of storage requirements is also a desirable feature of the matrix solution technique when the matrix is large and the computer storage capacity is limited.

Using matrix notation, Eq. 18 takes the form

$$AX = R \dots \dots \dots (24)$$

in which  $A$  = the coefficient matrix with components  $a_{ij}$  and  $X$ ,  $R$  are column vectors having components  $x_i$  and  $r_i$ , respectively, i.e.

$$A = \begin{bmatrix} a_{11} & a_{12} & & & & \\ a_{21} & a_{22} & a_{23} & a_{24} & & \\ a_{31} & a_{32} & a_{33} & a_{34} & & \\ & a_{43} & a_{44} & a_{45} & a_{46} & \\ & a_{53} & a_{54} & a_{55} & a_{56} & \\ \dots & \dots & \dots & \dots & \dots & \dots \\ & a_{2N-2,2N-3} & a_{2N-2,2N-2} & a_{2N-2,2N-1} & a_{2N-2,2N} & \\ & a_{2N-1,2N-3} & a_{2N-1,2N-2} & a_{2N-1,2N-1} & a_{2N-1,2N} & \\ & & & a_{2N,2N-1} & a_{2N,2N} & \end{bmatrix} \dots \dots (25)$$

$$\text{and } X = \begin{bmatrix} x_1 \\ x_2 \\ \cdot \\ \cdot \\ \cdot \\ x_{2N} \end{bmatrix} \quad \text{and } R = \begin{bmatrix} r_1 \\ r_2 \\ \cdot \\ \cdot \\ \cdot \\ r_{2N} \end{bmatrix} \dots \dots \dots (26)$$

If the components of  $A$  are shifted horizontally such that the relative positions of the components in any one row remain the same,  $A$  takes the form of  $A'$  with components  $a'_{ij}$ . Thus

$$A' = \begin{bmatrix} & & a'_{13} & a'_{14} \\ a'_{21} & a'_{22} & a'_{23} & a'_{24} \\ a'_{31} & a'_{32} & a'_{33} & a'_{34} \\ a'_{41} & a'_{42} & a'_{43} & a'_{44} \\ \vdots & \vdots & \vdots & \vdots \\ a'_2 N-2,1 & a'_2 N-2,2 & a'_2 N-2,3 & a'_2 N-2,4 \\ a'_2 N-1,1 & a'_2 N-1,2 & a'_2 N-1,3 & a'_2 N-1,4 \\ a'_2 N,1 & a'_2 N,2 & & \end{bmatrix} \dots \dots \dots (27)$$

This requires a relatively simple change of the  $j$ th index of the components of A.

The following technique efficiently solves the system of linear equations, denoted by Eq. 18, and now described by

$$A'X = R \dots \dots \dots (28)$$

The recurrent formulae, applicable to even-numbered rows, i.e. ( $i = 2, 4, 6, \dots 2N$ ) are

$$m_{i,2} = - a'_{i,1} \frac{m_{i-1,4}}{m_{i-1,3}} + a'_{i,2} \dots \dots \dots (29a)$$

$$z_i = - a'_{i,1} \frac{z_{i-1}}{m_{i-1,3}} + r_1 \dots \dots \dots (29b)$$

in which  $m_{1,3} = a'_{1,3}$ ,  $m_{1,4} = a'_{1,4}$  and  $z_1 = r_1$ . The recurrent formulae, applicable to the odd-numbered rows, i.e. ( $i = 3, 5, 7, \dots 2N - 1$ ) are:

$$m_{i,2} = - a'_{i,1} \frac{m_{i-2,4}}{m_{i-2,3}} + a'_{i,2} \dots \dots \dots (30a)$$

$$m_{i,3} = - a'_{i-1,3} \frac{m_{i,2}}{m_{i-1,2}} + a'_{i,3} \dots \dots \dots (30b)$$

$$m_{i,4} = - a'_{i-1,4} \frac{m_{i,2}}{m_{i-1,2}} + a'_{i,4} \dots \dots \dots (30c)$$

$$z_i = - m_{i,2} \frac{z_{i-1}}{m_{i-1,2}} - a'_{i,1} \frac{z_{i-2}}{m_{i-2,3}} + r_i \dots \dots \dots (30d)$$

The computations proceed sequentially from  $i = 2$  to  $i = 2N$ . The components of the solution vector, X, are obtained by back-substitution commencing at  $i = 2N$  and proceeding sequentially to  $i = 1$ . Thus

$$x_{2N} = \frac{z_{2N}}{m_{2N,2}} \dots \dots \dots (31)$$

and the recurrent formula for ( $i = 2N - 1, 2N - 3, \dots 5, 3, 1$ ) is

$$x_i = \frac{[z_i - m_{i,1} x_{i+1}]}{m_{i,3}} \dots \dots \dots (32)$$

while that for ( $i = 2N - 2, 2N - 4, \dots 6, 4, 2$ ) is

$$x_i = \frac{[z_i - a'_{i1} x_{i+2} - a'_{i2} x_{i+1}]}{m_{i2}} \dots\dots\dots (33)$$

When programing the preceding solution technique, it is not necessary to introduce the new components  $m_{ij}$  and  $z_i$  as these may be expressed as  $a'_{ij}$  and  $r_i$ , respectively.

*Acknowledgement.*

The work upon which this information is based was supported in part by funds provided under Project No. A-0-35-MO of the United States Department of the Interior, Office of Water Resources Research, as authorized under the Water Resources Research Act of 1964, for which T. E. Harbaugh was the principal project investigator.